Three grades of coherence for non-Archimedean preferences

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OUTLINE

- We review de Finetti's theory of coherence for a (weakly-ordered) preference relation ≤ over (bounded) real-valued random variables.
- de Finetti provides a representation of coherent₁, weak-ordered preferences using finitely additive, <u>real-valued</u> probabilities and expectations.
- Coherence₁ requires strict preference for a <u>uniformly dominating</u> variable. If for some $\varepsilon > 0$ and for each state ω , $X(\omega) + \varepsilon < Y(\omega)$ then X < Y.
 - 2. We examine two stronger coherence criteria: coherence₂ and coherence₃.
 - Coherence₂ requires strict preference for a <u>simply dominating</u> variable. If for each state ω , $X(\omega) < Y(\omega)$ then X < Y.
 - Coherence₃ requires strict preference for a <u>weakly dominating</u> variable (*admissibility*). If for each ω , $X(\omega) \leq Y(\omega)$, and for some ω , $X(\omega) < Y(\omega)$ then X < Y.

- We review how the two stronger coherence requirements (coherence₂ and coherence₃) impose unacceptable restrictions on <u>real-valued</u> preferences, as seen through the respective representations in terms of real-valued probabilities.
 - 3. We accommodate each of the three coherence conditions using (weakly ordered) non-Archimedean preferences.
- The central result is that each of these coherent weak-ordered preferences, over (even unbounded) real-valued random variables is represented using non-standard probability and utility.
- Simple examples illustrate that coherent₂ and coherent₃ weak orders cannot be represented using lexicographic probabilities and lexicographic utilities.
- Coherence₃ supports conditional probability and, more generally, conditional expectations derived entirely from unconditional preferences.
- Last, the same approach extends to represent non-Archimedean coherent *strict partial orders* yielding a version of non-standard IP theory.

1 – A short review of deFinetti's theory of coherent₁ wagering.

We have a zero-sum (sequential) game played between

a *Bookie* and a *Gambler*, with a *Moderator* supervising.

Let $X: \Omega \to \Re$ be a (bounded) real-valued variable defined on a space Ω of possibilities, a space that is well defined for all three players by the Moderator.

The *Bookie*'s <u>prevision</u> p(X) on the r.v. X has the operational content that, when the *Gambler* fixes a real-valued quantity $\Omega_{X, p(X)}$ then in state ω the resulting payoff to the *Bookie* is $\Omega_{X, p(X)} [X(\omega) - p(X)]$ with the opposite payoff to the Gambler.

• Given X, the *Bookie* offers a *fair price* (a constant variable) p(X) that makes the following two variables indifferent: $X \approx p(X)$

- A simple version of deFinetti's *Book* game proceeds as follows:
 - 1. The *Moderator* identifies a (possibly infinite) set of random variables $\{X_i\}$
 - 2. The *Bookie* announces a prevision $p_i = p(X_i)$ for each r.v. in the set.
 - 3. The *Gambler* then chooses (*finitely many*) non-zero terms $\alpha_i = \alpha_{X_i}, p(X_i)$.
 - 4. A state $\omega \in \Omega$ is realized by *Nature*: write $X_i(\omega) = X_i$
 - 5. The *Moderator* settles up each contract and awards the **Bookie** (Gambler) the respective SUM of his/her payoffs:

Total payoff to Bookie = $\sum_{i=1}^{n} \alpha_i [X_i - p_i]$.

Total payoff to Gambler = $-\sum_{i=1}^{n} \alpha_i [X_i - p_i]$.

Definition:

The *Bookie*'s previsions are <u>incoherent</u>₁ if the *Gambler* can choose <u>finitely</u> <u>many</u> non-zero terms, α_i that assures her/him a (*uniformly*) positive payoff, regardless which state in Ω obtains – and then the *Bookie* loses for sure. A set of previsions is <u>coherent</u>₁, if not incoherent₁.

Theorem (deFinetti):

A set of previsions is coherent₁ *if and only if* each prevision p(X) is the expectation for X under a common (finitely additive) probability *P*.

That is,
$$p(X) = E_{P(\bullet)}[X] = \int_{\Omega} X \, dP(\bullet)$$

• Corollary 1: When the random variables are *indicator functions* for events $\{E_i\}$, so that the gambles are simple bets – with the α 's then the stakes in a winner-take-all scheme – then

the previsions p_i are coherent₁ if and only if there is a (f.a.) probability P

where each prevision is the respective probability $p_i = P(E_i)$.

A short interlude

Three limitations to this general approach <u>not</u> addressed in this presentation

[1] *Dominance*(1, 2, or 3) is invalid in the presence of act-state dependence – *moral hazard*.
 Consider the following case of dominance(1, 2, or 3) between two acts.

Act A_2 dominates act A_1 . However, if there is *moral hazard* – act-state probabilistic dependence, then A_1 may maximize subjective conditional expected utility, not A_2 .

- Jiji Zhang's "Subjective Causal Networks ..." addresses this challenge!
- [2] <u>Strategic play</u> by the <u>Bookie</u> against the <u>Gambler</u> may result in a failure to elicit the <u>Bookie</u>'s degrees of belief. The <u>Bookie</u>'s prevision may differ from her/his credence.
- Kevin Zollman's "*The Theory of Games* ...," is one of several presentations at this workshop that intersect the theme: where strategic action conflicts with epistemic goals.
- [3] The problem of the <u>numeraire</u> state-dependent utilities.
 De Finetti's game is played with real-valued <u>unitless</u> outcomes of variables. The <u>Bookie</u>'s previsions may depend on which currency is used to realize variables.

- 2 What becomes of de Finetti's theory when
 - coherence₁ uniform dominance

is strengthened, either to

- coherence₂ strict dominance,
- or
- coherence₃ weak dominance (admissibility)?

Some Answers

• Coherence₂ precludes some intuitive (merely) f.a. expectations that are coherent₁. See de Finetti [1972, p. 77, *fn*.].

Example 1: Consider a denumerably infinite state space $\Omega = \{\omega_1, \omega_2, ... \}$.

Let P be the coherent₁ (strongly) finitely additive probability that is <u>uniform</u> on Ω .

For all integers *i* and *j*, $P(\{\omega_i\}) = P(\{\omega_j\})$. So, $P(\{\omega\}) = 0$.

Consider the (bounded) variable $X(\omega_n) = 1/n$. X is bounded: $0 < X \le 1$.

Then $\mathbf{E}_{P}[X] = 0$ and so, $X \approx 0$.

But for each ω , $X(\omega) > 0$. So X strictly dominates 0 and coherence₂ requires 0 < X.

Coherence₃ precludes a coherent₂ probability whenever some possible event is *P*-null.
 See Shimony [1955].

Example 2: Consider a binary space $\Omega = \{\omega_1, \omega_2\}$.

Let *P* be the coherent₂ probability supported by ω_1 : $P(\{\omega_1\}) = 1$ and $P(\{\omega_2\}) = 0$.

Let I_1 be the indicator for ω_1 : $I_1(\omega_1) = 1$ and $I_1(\omega_2) = 0$.

But, $E_P[I_1] = 1$ and then $I_1 \approx 1$.

However, I_1 is weakly dominated by the constant variable 1: $I_1(\omega_1) = 1$, $I_1(\omega_2) = 0$

 I_1 is *inadmissible* against the constant 1. Coherence₃ requires $I_1 < 1$.

How to accommodate all three coherence conditions without imposing such restrictive conditions on the probabilities that represent the corresponding uncertainties over Ω? 3. Non-Archimedean, weakly ordered coherent preference.

In de Finetti's theory of coherent previsions,

(independent of which of the three senses of "coherence" is applied) for each variable X, the *Bookie* is required to offer a <u>fair price</u> – a *prevision* p(X).

A *prevision* p(X) is a *real-valued constant* variable that, in the eyes of the *Bookie*, supports a swap, either way, between the variable X and the (constant) variable p(X).

• Expressed in terms of the *Bookie*'s preferences, $X \approx p(X)$.

Definition: A binary relation \leq is a *weak order* if it is transitive, and each pair of objects are comparable by the relation, i.e., either $X \leq Y$ or $Y \leq X$.

Definition: A binary relation \leq on a set \supseteq is a <u>Archimedean</u> – it admits a (real) Utility representation – if there exists a (real-valued) function $U: \supseteq \rightarrow \Re$ where, for $X, Y \in \supseteq$ $X \leq Y$ if and only if $U(X) \leq U(Y)$.

Given a real number *c*, let $C(\omega) = c$ denote the corresponding constant variable. Under each of the three dominance conditions, when c < d, then C < D.

Let \leq be a *coherent*, weakly ordered preference over (bounded) variables defined on a common state-space Ω . Because of the requirements for trading, under coherence:

If for each variable X there is a real-valued prevision, p(X), then ≤ is Archimedean with U(X) = p(X),
 and for real α and β U(αX + βY) = αU(X) + βU(Y) = αp(X) + βp(Y).

We modify de Finetti's prevision-game by *not requiring* that the *Bookie* holds a real-valued (constant) prevision, p(X), for swapping with variable X.

• Aside: We do not require that the variables are bounded.

Instead, we require that the *Bookie* has a coherent_(1,2, or 3) weakly ordered preference \leq over the set of real-valued variables defined on a state-space Ω , and where the *Bookie* accepts permitted trades just as in de Finetti's *Prevision Game*.

Coherence_(1,2, or 3) of \leq means that the corresponding dominance_(1, 2, or 3) condition is respected and that the following *Independence* condition is satisfied.

Let *a* and *b* be real numbers, *c* a positive real number, and *Y* a variable.

 $X_1 \approx X_2$ if and only if $aX_1 + bY \approx aX_2 + bY$ $X_1 \prec X_2$ if and only if $cX_1 + bY \prec cX_2 + bY$ The weak order \leq is operationalized through allowed trades,

with the *Bookie* specifying:

(1) the strict partial order $X \prec Y$, which identifies all the 1-way trades.

Where the *Bookie* is willing to swap X for Y, but not vice versa.

(2) the equivalence relation, $X \approx Y$, which identifies all the 2-way trades:

Where the *Bookie* is willing to swap X for Y, and willing to swap Y for X.

*Two illustrations of non-Archimedean, coherent*₃ *weak orders*

Example 3: Let $\Omega = \{\omega_1, \omega_2\}$. Variable X_i is the ordered pair $\langle x_{i1}, x_{i2} \rangle$, where $X_i(\omega_j) = x_{ij}$.

Define the coherent³ ≤ weak order where

(1) $X_1 \prec X_2$ iff $x_{11} + x_{12} < x_{21} + x_{22}$

or $x_{11} + x_{12} = x_{21} + x_{22}$ and $x_{11} < x_{21}$.

(2) $X_1 \approx X_2$ iff $X_1 = X_2$



Increasing preference for points to the SE on a given line.



 ω_2

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Example 4 (de Finetti's *Example 1* continued):

Consider, again, the denumerable state space $\Omega = \{\omega_1, \omega_2, \ldots\}$.

A coherent₃ preference allows indifference between state-indicators, $I(\omega_i) \approx I(\omega_j)$.

Each is strictly more desirable than 0 (by dominance₃), $0 \prec I(\omega_i)$.

And by dominance₂, each is strictly less desirable than $X(\omega_n) = 1/n$, $I(\omega_j) \prec X$,

By coherence₃, X is less desirable than an arbitrary positive constant $C(\omega) = c > 0$.

$$0 \prec I(\omega_i) \approx I(\omega_j) \prec X \prec C.$$

The following shows that the two weak orders of Examples 3 and 4 are non-Archimedean.

Definition: Let ▷ be a total order on a set \aleph . A subset $\beth \subseteq \aleph$ is <u>▷-order dense</u> if, whenever $x \triangleright y$ with $x \notin \beth$ and $y \notin \beth$, then there exists $z \in \beth$ with $x \triangleright z \triangleright y$.

Lemma (Fishburn, 1972):

Let \triangleright be a total order on a set \aleph . \triangleright is an Archimedean order (i.e., \triangleright has a real-valued Utility representation) *if and only if* there is a denumerable \triangleright -order dense subset $\supseteq \subseteq \aleph$.

Recall: When \leq is a weak order, then $\triangleright = \leq /\approx$ is a total order.

Next we see that there are no denumerable order dense subsets either in Example 3 or 4. The reasoning is similar in both examples. Let \supseteq be an \triangleright -order dense subset and consider the continuum-many reals, 0 < c < 1.

In *Example* 3, each *c* denotes a distinct set (a line with slope -1) of continuum-many variables X_{α}^{c} where $x_{\alpha,1}^{c} + x_{\alpha,2}^{c} = c$. Since $X_{\alpha}^{c} \prec X_{b}^{c}$ whenever $x_{\alpha,1}^{c} \prec x_{b,1}^{c}$, if \supseteq is an \triangleright -order dense subset, then \supseteq contains at least one point from line *c*. Hence \supseteq is uncountable.

In *Example* 4, for each pair of real numbers $0 \le c \le d \le 1$, the \le order satisfies

$$0 < C < C + I(\omega_1) < D < D + I(\omega_1).$$

If \supseteq is an \triangleright -order dense subset, then for each 0 < c < 1, \supseteq contains at least one variable, X_c , where $C \leq X_c \leq C + I(\omega_1)$. Hence \supseteq is uncountable. Representing non-Archimedean, coherent weak-orders using non-standard utilities

Let Ξ be a linear space of (standard) real-valued variables on a set Ω .

Let X and Y belong to Ξ , and let a and b be (standard) real numbers.

Denote the non-standard real numbers by *R.

Aside: We use the ultra-product model of *R.

Fix one of the three senses of *dominance*.

Defn: A non-standard value function $U: \Xi \to \mathscr{R}$ is a *positive linear functional* if

- whenever *Y* dominates *X* then U(X) < U(Y)
- U(aX + bY) = aU(X) + bU(Y)
- *Main Theorem*: Let \leq be a weak order over Ξ .

 \leq is coherent_(1, 2, or 3) if and only if

there is a positive linear functional U that represents \leq .

Non-standard Probability defined by non-standard Utility.

In the *Main Theorem*, without loss of generality, one may "standardize" the non-standard utility U so that U(0) = 0 and U(1) = 1.

Then, as *U* is a positive linear functional, when restricted to indicator variables, that is, with I_E the indicator function for the event $E (\subseteq \Omega)$ then

• U is a (finitely additive) non-standard probability, U = *P.

Let *E* and *F* be disjoint subsets of Ω , with $G = E \cup F$. Then *U* satisfies:

> $U(I_{\Omega}) = 1$ and $U(I_{\emptyset}) = 0$ $0 \leq U(I_E) \leq 1$ $U(I_G) = U(I_E) + U(I_F)$

Definition: Call a non-standard ε a *positive infinitesimal* if $0 < \varepsilon < a$ for each positive (standard) real number a.

Example 2 (continued):

Let $\Omega = \{\omega_1, \omega_2\}$. Variable X_i is the ordered pair $\langle x_{i1}, x_{i2} \rangle$, where $X_i(\omega_j) = x_{ij}$. Define the coherent₃ \leq weak order where

- (1) $X_1 < X_2$ iff $x_{11} + x_{12} < x_{21} + x_{22}$ or $x_{11} + x_{12} = x_{21} + x_{22}$ and $x_{11} < x_{21}$.
- (2) $X_1 \approx X_2$ iff $X_1 = X_2$.

Increasing preference for all points on lines to the NE across different lines.

Increasing preference for points to the SE on a given



The representing non-standard *P-probability satisfies

$$*P(I_{\omega_1}) = \frac{1}{2} + \varepsilon$$
, and $*P(I_{\omega_2}) = \frac{1}{2} - \varepsilon$

for a positive infinitesimal ε.

Some coherent preferences that have no lexicographic-probability representations.

Let $\vec{P} = \langle P_1, P_2, ... \rangle$ be a well-ordered sequence of real-valued probabilities. Each probability P_i is defined on an algebra \mathcal{A} of subsets of a state-space Ω . For events *E*, *F* that belong to \mathcal{A} , say that

 $E \Join_{\overrightarrow{P}} F \text{ iff } P_i(E) < P_i(F) \text{ for the least index } i \text{ where } P_i(E) \neq P_i(F)$ and $E \approx_{\overrightarrow{P}} F \text{ iff for all } i, P_i(E) = P_i(F).$

Write $E \simeq_{\overrightarrow{P}} F$ to abbreviate $E \simeq_{\overrightarrow{P}} F$ or $E \approx_{\overrightarrow{P}} F$.

• These relations define a qualitative (non-Archimedean) probability.

- (1) $\bowtie_{\vec{P}}$ is a weak-order
- (2) $0 \cong_{\overrightarrow{P}} E \cong_{\overrightarrow{P}} 1$
- (3) $E \simeq_{\overrightarrow{P}} F$ if and only if $E \cup G \simeq_{\overrightarrow{P}} F \cup G$ whenever $E \cap G = F \cap G = \emptyset$.

Example 5: Let Ω be an infinite set. Consider a coherent₃ weak-order \leq where $I_{\omega_{\alpha}} \approx I_{\omega_{\beta}}$ for each pair of states in Ω , [1] and (by *admissibility*) $I_{\varnothing} < I_{\omega_{\alpha}}$. [2] Necessary for a lex-prob \vec{P} to agree with [1], $P_i(\{\omega\}) = 0$ for all i and all $\omega \in \Omega$. But then $I_{\varnothing} \approx_{\vec{P}} I_{\omega_{\alpha}}$, contrary to what [2] requires.

Also, there are coherent₂ weak-orders which are not coherent₃ and which are not represented by a Utility based on a lex-probability.

Example 6 (de Finetti's *Example* 1, modified) Consider the denumerable $\Omega = \{\omega_1, \omega_2, ...\}$. A coherent₂ preference allows indifference between each two state-indicators, $I(\omega_m)$ and $I(\omega_n)$, and also indifference between each state indicator and 0,

$$\mathbf{0} \approx I(\omega_m) \approx I(\omega_n)$$
.

But by coherence₂ (by strict dominance) each of these variables is strictly less desirable than $X(\omega_n) = 1/n$, which yields

$$\mathbf{0} \approx I(\omega_m) \approx I(\omega_n) \prec X.$$

Necessary for a lex-prob \vec{P} to agree with these indifferences is that $P_i(\{\omega_m\}) = 0$ for all *i* and all *m*. But then for each *i*, the expectation $E_{P_i}[X] = 0$, and so $0 \approx_{\vec{P}} X$, contrary to strict dominance (coherence₂) which requires 0 < X.

On conditional expectations:

Reconsider de Finetti's Prevision Game

Definition: A called-off prevision p(X || E) for X, made by the *Bookie* has a payoff scheme to the Bookie: $\alpha_X ||_E E(\omega) [X(\omega) - p(X || E)].$

Corollary 2 to de Finetti's Coherence Theorem:

A called-off prevision p(X || E) is coherent₁ alongside the (coherent₁) previsions p(X) for X, and when p(E) > 0, *if and only if*

 $p(X \parallel E)$ is the *conditional expectation* under *P* for *X*, given *E*.

That is, $p(X || E) = E_{P(\bullet | E)}[X] = \int_{\Omega} X dP(\bullet | E)$.

But, when p(E) = 0, p(X || E) is <u>unconstrained</u> by coherence₁ or coherence₂.

The situation is different with coherence₃.

Coherence₃ entails that if $E \neq \emptyset$ then 0 < E. Using the non-standard representation of the *Main Theorem*: If $E \neq \emptyset$ then 0 < *P(F).

Then, just as in the real-valued theory given an event *E* of positive probability, *called-off* preference fixes conditional *probability, given *E*.

That is, $*P(F | E) = *P(F \cap E) / *P(E)$.

An introduction to *Imprecise non-standard Probability*: *I*P*-theory.

For a final theme, consider another de Finetti result, which addresses extending coherent₁ previsions to a larger collection of variables.

<u>Background</u>: The <u>Bookie</u> assigns coherent₁ previsions to a set χ of (bounded) variables. For each $X \in \chi$ the <u>Bookie</u> assigns a prevision p(X), and these are coherent₁. By the rules of the <u>Prevision Game</u>, these previsions determine a unique coherent₁ previsions for each variable Y in the Linear Span[χ].

Let Z be a (bounded) variable defined on Ω but outside the *Linear Span*[χ].

Define: $\underline{Z} = \{X: X(\omega) \le Z(\omega) \text{ and } X \text{ in the } Linear Span[\chi]\} - \text{ to approximate from below.}$ $\overline{Z} = \{X: X(\omega) \ge Z(\omega) \text{ and } X \text{ in the } Linear Span[\chi] - \text{ to approximate from above.}$ Let $\underline{p}(Z) = \sup_{X \in \underline{Z}} p(X) \text{ and } \overline{p}(Z) = \inf_{X \in \overline{Z}} p(X)$

Fundamental Theorem of Previsions

To remain coherent₁, p(Z) may be any value in the closed interval $[\underline{p}(Z), \overline{p}(Z)].$

Adapting the Fundamental Theorem to I*P theory using a coherent pre-order.

A *<u>coherent pre-order</u>* is operationalized by a partition into 4 categories of trades:

- a strict partial order X < Y, which identifies all the 1-way trades.
 Where the *Bookie* is willing to swap X for Y, but not vice versa.
- (2) an equivalence relation, $X \approx Y$, which identifies all the 2-way trades: Where the *Bookie* is willing to swap X for Y, and willing to swap Y for X.
- (3) the <u>one-way limited</u> trades, X ≥ Y,
 This identifies those pairs where the <u>Bookie</u> is willing to swap one way,
 Y for X, but has not resolved whether to trade the other way, X for Y.
- (4) the non-comparable pairs, $X \Leftrightarrow Y$, where the *Bookie* is unwilling to trade either way.

Note: When both categories (3) and (4) are empty, the pre-order is a weak-order.

In order to be *coherent*_(1, 2, or 3) the pre-order respects the (respective) dominance condition and satisfies the *Independence* condition for allowed trades.

Let a and b be real numbers, c a positive real number, and Y a variable.

 $X_1 \approx X_2$ if and only if $aX_1 + bY \approx aX_2 + bY$ $X_1 \prec X_2$ if and only if $cX_1 + bY \prec cX_2 + bY$ A *coherent extension* of a pre-order preserves all the binary comparisons already fixed by categories (1) and (2), and either moves some comparisons from category (3) into category (1) or (2), or moves some comparisons from category (4) into category (1), (2), or (3) while satisfying the (respective) *Coherence* and *Independence* conditions.

We show how to represent a coherent pre-order by the set of all its coherent, weak-order extensions – the extensions where categories (3) and (4) are empty.

Then, by the *Main Theorem*, each coherent pre-order is represented by the set of non-standard utilities *U* that correspond to each coherent weak-order extension.

• Each coherent pre-order is represented as an I*P set of non-standard utilities.

<u>Summary</u>

- We generalize de Finetti's theory of
 - coherent₁, real-valued previsions over (bounded) real-valued variables
- to a theory of

coherent_(1,2, or 3) non-Archimedean, weak-orders over real-valued variables.

- 1 This theory accommodates stronger dominance conditions without having to impose overly restrictive conditions on personal probabilities.
- 2 A non-standard Utility represents a coherent_(1,2, or 3) non-Archimedean weak-order. The Utility reduces to a non-standard probability over events (indicator variables).
- 3 Not all non-Archimedean weak orders are representable using lex-probabilities.
- 4 With coherence₃, all conditional probabilities are fixed by unconditional probabilities.
- 5 We adapt de Finetti's *Fundamental Theorem* to apply to coherent pre-orders, which opens the door to *Imprecise non-standard Probability theory*: I*P- theory